

# KADM5 Library and Server Implementation Design\*

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# 1 Overview

The KADM5 administration system is designed around the KADM5 API. The “server-side” library `libkadm5srv.a` implements the KADM5 API by operating directly on the underlying KDC and admin databases. The “client-side” library `libkadm5clnt.a` implements the KADM5 API via an RPC mechanism. The administration server `kadmind` accepts RPC requests from the client-side library and translates them into calls to the server-side library, performing authentication, authorization, and logging along the way.

The two libraries, `libkadm5clnt.a` and `libkadm5srv.a`, export the identical `kadm5` interface; for example, both contain definitions for `kadm5_get_principal`, and all other `kadm5` functions. In most cases, the client library function just marshalls arguments and results into and out of an RPC call, whereas the server library function performs the actual operation on the database file. `kadm5_init_*`, however, are substantially different even though they export the same interface: on the client, they establish the RPC connection and GSS-API context, whereas on the server side they open the database files, read in the password dictionary, and the like. Also, the `kadm5_free` functions operate on local process memory in both libraries.

The admin server is implemented as a nearly-stateless transaction server, where each admin API function represents a single transaction. No per-client or per-connection information is stored; only local database handles are maintained between requests. The RPC mechanism provides access to remote callers’ authentication credentials for authorization purposes.

The admin API is exported via an RPC interface that hides all details about network encoding, authentication, and encryption of data on the wire. The RPC mechanism does, however,

allow the server to access the underlying authentication credentials for authorization purposes.

The admin system maintains two databases:

- The master Kerberos (KDC) database is used to store all the information that the Kerberos server understands, thus allowing the greatest functionality with no modifications to a standard KDC.
- The KDC database also stores kadm5-specific per-principal information in each principal's `krb5_tl_data` list. In a prior version, this data was stored in a separate admin principal database; thus, when this document refers to “the admin principal database,” it now refers to the appropriate `krb5_tl_data` entries in the KDC database.
- The policy database stores kadm5 policy information.

The per-principal information stored in the admin principal database consists of the principal's policy name and an array of the principal's previous keys. The old keys are stored encrypted in the key of the special principal “`kadmin/history`” that is created by the server library when it is first needed. Since a change in `kadmin/history`'s key renders every principal's key history array useless, it can only be changed using the `ovsec_adm_edit` utility; that program will reencrypt every principal's key history in the new key.<sup>1</sup> The server library refuses all requests to change `kadmin/history`'s key.

## 2 API Handles

Each call to `kadm5_init_*` on the client or server creates a new API handle. The handles encapsulate the API and structure versions specified by `kadm5_init_*`'s caller and all other internal data needed by the library. A process can have multiple open API handles simultaneously by calling `kadm5_init_*` multiple times, and call can specify a different version, client or service principal, and so forth.

Each `kadm5` function verifies the handle it is given with the `CHECK_HANDLE` or `_KADM5_CHECK_HANDLE` macros. The `CHECK_HANDLE` macro differs for the client and server library because the handle types used by those libraries differ, so it is defined in both `<client.internal.h>` and `<server.internal.h>` in the library source directory. In each header file, `CHECK_HANDLE` first calls `GENERIC_CHECK_HANDLE`, defined

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<sup>1</sup>`ovsec_adm_edit` has not yet been implemented, and there are currently no plans to implement it; thus, the history cannot currently be changed.

in `<admin_internal.h>`, which verifies the magic number, API version, and structure version that is contained in both client and server handles. `CHECK_HANDLE` then calls either `CLIENT_CHECK_HANDLE` or `SERVER_CHECK_HANDLE` respectively to verify the client- or server-library specific handle fields.

The `CHECK_HANDLE` macro is useful because it inlines the handle check instead of requiring a separate function call. However, using `CHECK_HANDLE` means that a source file cannot be compiled once and included into both the client and server library, because `CHECK_HANDLE` is always either specific to either the client or server library, not both. There are a number of functions that can be implemented with the same code in both the client and server libraries, however, including all of the `kadm5_free` functions and `kadm5_chpass_principal_util`. The `_KADM5_CHECK_HANDLE` macro solves this problem; instead of inlining the handle check, it calls the function `_kadm5_check_handle` which is defined separately in both the client and server library, in `client_init.c` and `server_init.c`. Since these two files are only compiled once and put in a single library, they simply verify the handle they are passed with `CHECK_HANDLE` and return the result.

### 3 API Versioning

The KADM5 system is designed to support multiple versions of the KADM5 API. Presently, two versions exist: `KADM5_API_VERSION_1` and `KADM5_API_VERSION_2`. The former is equivalent to the initial OpenVision API, `OVSEC_KADM_API_VERSION_1`; the latter was created during the initial integration of the OpenVision system into the MIT release.

Implementing a versioned API in C via with both local and RPC access presents a number of design issues, some of them quite subtle. The contexts in which versioning considerations must be made include:

1. Typedefs, function declarations, and defined constants depend on the API version a client is written to and must be correct at compile time.
2. Each function in the server library must behave according to the API version specified by the caller at runtime to `kadm5_init_*`.
3. The XDR functions used by the RPC layer to transmit function arguments and results must encode data structures correctly depending on the API version specified by the client at runtime.
4. Each function in the client library must behave according to the API version specified by the caller at runtime to `kadm5_init_*`.

5. The RPC server (kadmind) must accept calls from a client using any supported API version, and must then invoke the function in the server library corresponding to the RPC with the API version indicated by the client caller.
6. When a first API function is invoked that needs to call a second function in the API on its own behalf, and that second API function's behavior depends on the API version specified, the first API function must either be prepared to call the second API function at whatever version its caller specifies or have a means of always calling the second API function at a pre-determined version.

The following functions describe how each context is handled.

### 3.1 Designing for future compatibility

Any code whose behavior depends on the API version should be written so as to be compatible with future, currently unknown API versions on the grounds that any particularly piece of API behavior will most likely not change between versions. For example, in the current system, the code is not written as “if this is VERSION\_1, do X, else if this is VERSION\_2, do Y”; instead, it is written as “if this is VERSION\_1, do X; else, do Y.” The former will require additional work when VERSION\_3 is defined, even if “do Y” is still the correct action, whereas the latter will work without modification in that case.

### 3.2 Header file declarations

Typedefs, defined constants and macros, and function declarations may change between versions. A client is always written to a single, specific API version, and thus expects the header files to define everything according to that API. Failure of a header file to define values correctly will result in either compiler warnings (e.g. if the pointer type of a function argument changes) or fatal errors (e.g. if the number of arguments to a function changes, or the fields of a structure change). For example, in VERSION\_1, `kadm5_get_policy` took a pointer to a pointer to a structure, and in VERSION\_2 it takes a pointer to a structure; that would generate a warning if not correct. In VERSION\_1, `kadm5_randkey_principal` accepted three arguments but in VERSION\_2 accepts four; that would generate a fatal error.

The header file defines everything correctly based on the value of the `USE_KADM5_API_VERSION` constant. The constant can be assigned to an integer corresponding to any supported API version, and defaults to the newest version. The header files then simply use an `#ifdef` to include the right definitions:

```
#if USE_KADM5_API_VERSION == 1
```

```

kadm5_ret_t    kadm5_get_principal(void *server_handle,
                                   krb5_principal principal,
                                   kadm5_principal_ent_t *ent);

#else
kadm5_ret_t    kadm5_get_principal(void *server_handle,
                                   krb5_principal principal,
                                   kadm5_principal_ent_t ent,
                                   long mask);

#endif

```

### 3.3 Server library functions

Server library functions must know how many and what type of arguments to expect, and must operate on those arguments correctly, based on the API version with which they are invoked. The API version is contained in the handle that is always passed as their first argument, generated by `kadm5_init_*` (to which the client specified the API version to use at run-time).

In general, it is probably unsafe for a compiled function in a library to re-interpret the number and type of defined arguments at run-time since the calling conventions may not allow it; for example, a function whose first argument was a short in one version and a pointer in the next might fail if it simply typed-casted the argument. In that case, the function would have to be written to take variable arguments (i.e. use `<stdarg.h>`) and extract them from the stack based on the API version. Alternatively, a separate function for each API version could be defined, and `<kadm5/admin.h>` could be written to `#define` the exported function name based on the value of `USE_KADM5_API_VERSION`.

In the current system, it turns out, that isn't necessary, and future implementors should take care to ensure that no version has semantics that will cause such problems in the future. All the functions in KADM5 that have different arguments or results between `VERSION_1` and `VERSION_2` do so simply by type-casting their arguments to the appropriate version and then have separate code paths to handle each one correctly. `kadm5_get_principal`, in `svr_principal.c`, is a good example. In `VERSION_1`, it took the address of a pointer to a `kadm5_principal_ent_t` to fill in with a pointer to allocated memory; in `VERSION_2`, it takes a pointer to a structure to fill in, and a mask of which fields in that structure should be filled in. Also, the contents of the `kadm5_principal_ent_t` changed slightly between the two versions. `kadm5_get_principal` handles versioning as follows (following along in the source code will be helpful):

1. If `VERSION_1`, it saves away its entry argument (address of a pointer to a structure) and resets its value to contain the address of a locally stack-allocated entry structure;

this allows most of the function to be written once, in terms of `VERSION_2` semantics. If `VERSION_1`, it also resets its mask argument to be `KADM5_PRINCIPAL_NORMAL_MASK`, because that is the equivalent to `VERSION_1` behavior, which was to return all the fields of the structure.

2. The bulk of the function is implemented as expected for `VERSION_2`.
3. The new fields in the `VERSION_2` entry structure are assigned inside a block that is only executed if the caller specified `VERSION_2`. This saves a little time for a `VERSION_1` caller.
4. After the entry structure is filled, the function checks again if it was called as `VERSION_1`. If so, it allocates a new `kadm5_principal_ent_t_v1` structure (which is conveniently defined in the header file) with `malloc`, copies the appropriate values from the entry structure into the `VERSION_1` entry structure, and then writes the address of the newly allocated memory into address specified by the original entry argument which it had previously saved away.

There is another complication involved in a function re-interpreting the number of arguments it receives at compile time—it cannot assign any value to an argument for which the client did not pass a value. For example, a `VERSION_1` client only passes three arguments to `kadm5_get_principal`. If the implementation of `kadm5_get_principal` notices that the caller is `VERSION_1` and therefore assigns its fourth argument, `mask`, to a value that mimics the `VERSION_1` behavior, it may inadvertently overwrite data on its caller's stack. This problem can be avoided simply by using a true local variable in such cases, instead of treating an unpassed argument as a local variable.

### 3.4 XDR functions

The XDR functions used to encode function arguments and results must know how to encode the data for any API version. This is important both so that all the data gets correctly transmitted and so that protocol compatibility between clients or servers using the new library but an old API version is maintained; specifically, new `kadmind` servers should support old `kadm5` clients.

The signature of all XDR functions is strictly defined: they take the address of an XDR function and the address of the data object to be encoded or decoded. It is thus impossible to provide the API version of the data object as an additional argument to an XDR function. There are two other means to convey the information, storing the API version to use as a field in the data object itself and creating separate XDR functions to handle each different version of the data object, and both of them are used in `KADM5`.

In the client library, each `kadm5` function collects its arguments into a single structure to be passed by the RPC; similarly, it expects all of the results to come back as a single structure from the RPC that it will then decode back into its constituent pieces (these are the standard ONC RPC semantics). In order to pass versioning information to the XDR functions, each function argument and result datatype has a field to store the API version. For example, consider `kadm5_get_principal`'s structures:

```
struct gprinc_arg {
    krb5_ui_4 api_version;
    krb5_principal princ;
    long mask;
};
typedef struct gprinc_arg gprinc_arg;
bool_t xdr_gprinc_arg();

struct gprinc_ret {
    krb5_ui_4 api_version;
    kadm5_ret_t code;
    kadm5_principal_ent_rec rec;
};
typedef struct gprinc_ret gprinc_ret;
bool_t xdr_gprinc_ret();
```

`kadm5_get_principal` (in `client_principal.c`) assigns the `api_version` field of the `gprinc_arg` to the version specified by its caller, assigns the `princ` field based on its arguments, and assigns the `mask` field from its argument if the caller specified `VERSION_2`. It then calls the RPC function `clnt_call`, specifying the XDR functions `xdr_gprinc_arg` and `xdr_gprinc_ret` to handle the arguments and results.

`xdr_gprinc_arg` is invoked with a pointer to the `gprinc_arg` structure just described. It first encodes the `api_version` field; this allows the server to know what to expect. It then encodes the `krb5_principal` structure and, if `api_version` is `VERSION_2`, the `mask`. If `api_version` is not `VERSION_2`, it does not encode *anything* in place of the `mask`, because an old `VERSION_1` server will not expect any other data to arrive on the wire there.

The server performs the `kadm5_get_principal` call and returns its results in an XDR encoded `gprinc_ret` structure. `clnt_call`, which has been blocking until the results arrived, invokes `xdr_gprinc_ret` with a pointer to the encoded data for it to decode. `xdr_gprinc_ret` first decodes the `api_version` field, and then the `code` field since that is present in all versions to date. The `kadm5_principal_ent_rec` presents a problem, however. The structure does not itself contain an `api_version` field, but the structure is different between the two versions. Thus, a single XDR function cannot decode both versions of the structure because it will

have no way to decide which version to expect. The solution is to have two functions, `kadm5_principal_ent_rec_v1` and `kadm5_principal_ent_rec`, which always decode according to `VERSION_1` or `VERSION_2`, respectively. `gprinc_ret` knows which one to invoke because it has the `api_version` field returned by the server (which is always the same as that specified by the client in the `gpring_arg`).

In hindsight, it probably would have been better to encode the API version of all structures directly in a version field in the structure itself; then multiple XDR functions for a single data type wouldn't be necessary, and the data objects would stand complete on their own. This can be added in a future API version if desired.

### 3.5 Client library functions

Just as with server library functions, client library functions must be able to interpret their arguments and provide result according to the API version specified by the caller. Again, `kadm5_get_principal` (in `client_principal.c`) is a good example. The `gprinc_ret` structure that it gets back from `clnt_call` contains a `kadm5_principal_ent_rec` or a `kadm5_principal_ent_rec_v1` (the logic is simplified somewhat because the `VERSION_2` structure only has new fields added on the end). If `kadm5_get_principal` was invoked with `VERSION_2`, that structure should be copied into the pointer provided as the entry argument; if it was invoked with `VERSION_1`, however, the structure should be copied into allocated memory whose address is then written into the pointer provided by the entry argument. Client library functions make this determination based on the API version specified in the provided handle, just like server library functions do.

### 3.6 Admin server stubs

When an RPC call arrives at the server, the RPC layer authenticates the call using the GSS-API, decodes the arguments into their single-structure form (ie: a `gprinc_arg`) and dispatches the call to a stub function in the server (in `server_stubs.c`). The stub function first checks the caller's authorization to invoke the function and, if authorized, calls the `kadm5` function corresponding to the RPC function with the arguments specified in the single-structure argument.

Once again, `kadm5_get_principal` is a good example for the issues involved. The contents of the `gprinc_arg` given to the stub (`get_principal_1`) depends on the API version the caller on the client side specified; that version is available to the server in the `api_version` field of the `gprinc_arg`. When the server calls `kadm5_get_principal` in the server library, it must give that function an API handle that contains the API version requested by the client; otherwise the function semantics might not be correct. One possibility would be for the

server to call `kadm5_init` for each client request, specifying the client's API version number and thus generating an API handle with the correct version, but that would be prohibitively inefficient. Instead, the server dips down in the server library's internal abstraction barrier, using the function `new_server_handle` to cons up a server handle based on the server's own `global_server_handle` but using the API version specified by the client. The server then passes the newly generated handle to `kadm5_get_principal`, ensuring the right behavior, and creates the `gprinc_ret` structure in a manner similar to that described above.

Although `new_server_handle` solves the problem of providing the server with an API handle containing the right API version number, it does not solve another problem: that a single source file, `server_stubs.c`, needs to be able to invoke functions with arguments appropriate for multiple API versions. If the client specifies `VERSION_1`, for example, the server must invoke `kadm5_get_principal` with three arguments, but if the client specifies `VERSION_2` the server must invoke `kadm5_get_principal` with four arguments. The compiler will not allow this inconsistency. The server defines wrapper functions in a separate source file that match the old version, and the separate source file is compiled with `USE_KADM5_API_VERSION` set to the old version; see `kadm5_get_principal_v1` in `server_glue_v1.c`. The server then calls the correct variant of `kadm5_get_principal_*` based on the API version and puts the return values into the `gprinc_ret` in a manner similar to that described above.

Neither of these solutions are necessarily correct. `new_server_handle` violates the server library's abstraction barrier and is at best a kludge; the server library should probably export a function to provide this behavior without violating the abstraction; alternatively, the library should be modified so that having the server call `kadm5_init` for each client RPC request would not be too inefficient. The glue functions in `server_glue_v1.c` really are not necessary, because the server stubs could always just pass dummy arguments for the extra arguments; after all, the glue functions pass *nothing* for the extra arguments, so they just end up as stack garbage anyway.

Another alternative to the `new_server_handle` problem is to have the server always invoke server library functions at a single API version, and then have the stubs take care of converting the function arguments and results back into the form expected by the caller. In general, however, this might require the stubs to duplicate substantial logic already present in the server library and further violate the server library's abstraction barrier.

### 3.7 KADM5 self-reference

Some `kadm5` functions call other `kadm5` functions “on their own behalf” to perform functionality that is necessary but that does not directly affect what the client sees. For example, `kadm5_chpass_principal` has to enforce password policies; thus, it needs to call `kadm5_get_principal` and, if the principal has a policy, `kadm5_get_policy` and `kadm5_modify_principal` in the process of changing a principal's password. This leads to a complication: what API

handle should `kadm5_chpass_principal` pass to the other `kadm5` functions it calls?

The “obvious,” but wrong, answer is that it should pass the handle it was given by its caller. The caller may provide an API handle specifying any valid API version. Although the semantics of `kadm5_chpass_principal` did not change between `VERSION_1` and `VERSION_2`, the declarations of both `kadm5_get_principal` and `kadm5_get_policy` did. Thus, to use the caller’s API handle, `kadm5_chpass_principal` will have to have a separate code path for each API version, even though it itself did not change between versions, and duplicate a lot of logic found elsewhere in the library.

Instead, each API handle contains a “local-use handle,” or `lhandle`, that `kadm5` functions should use to call other `kadm5` functions. For example, the client-side library’s handle structure is:

```
typedef struct _kadm5_server_handle_t {
    krb5_ui_4      magic_number;
    krb5_ui_4      struct_version;
    krb5_ui_4      api_version;
    char *         cache_name;
    int            destroy_cache;
    CLIENT *       clnt;
    krb5_context    context;
    kadm5_config_params params;
    struct _kadm5_server_handle_t *lhandle;
} kadm5_server_handle_rec, *kadm5_server_handle_t;
```

The `lhandle` field is allocated automatically when the handle is created. All of the fields of the API handle that are accessed outside `kadm5_init` are also duplicated in the `lhandle`; however, the `api_version` field of the `lhandle` is always set to a *constant* value, regardless of the API version specified by the caller to `kadm5_init`. In the current implementation, the `lhandle`’s `api_version` is always `VERSION_2`.

By passing the caller’s handle’s `lhandle` to recursively called `kadm5` functions, a `kadm5` function is assured of invoking the second `kadm5` function with a known API version. Additionally, the `lhandle`’s `lhandle` field points back to the `lhandle`, in case `kadm5` functions call themselves more than one level deep; `handle->lhandle` always points to the same `lhandle`, no matter how many times the indirection is performed.

This scheme might break down if a `kadm5` function has to call another `kadm5` function to perform operations that they client will see and for its own benefit, since the semantics of the recursively-called `kadm5` function may depend on the API version specified and the client may be depending on a particular version’s behavior. Future implementors should avoid creating a situation in which this is possible.

## 4 Server Main

The admin server starts by trapping all fatal signals and directing them to a cleanup-and-exit function. It then creates and exports the RPC interface and enters its main loop.

The main loop dispatches all incoming requests to the RPC mechanism. In a previous version, after 15 seconds of inactivity, the server closed all open databases; each database was be automatically reopened by the API function implementations as necessary. That behavior existed to protect against loss of written data before the process exited. The current database libraries write all changes out to disk immediately, however, so this behavior is no longer required or performed.

## 5 Remote Procedure Calls

The RPC for the Admin system will be based on ONC RPC. ONC RPC is used because it is a well-known, portable RPC mechanism. The underlying external data representation (xdr) mechanisms for wire encapsulation are well-known and extensible. Authentication to the admin server and encryption of all RPC functional arguments and results are be handled via the AUTH\_GSSAPI authentication flavor of ONC RPC.

## 6 Database Record Types

### 6.1 Admin Principal, `osa_princ_ent_t`

The admin principal database stores records of the type `osa_princ_ent_t` (declared in `<kadm5/adb.h>`), which is the subset of the `kadm5_principal_ent_t` structure that is not stored in the Kerberos database plus the necessary bookkeeping information. The records are keyed by the ASCII representation of the principal's name, including the trailing NULL.

```
typedef struct _osa_pw_hist_t {
    int n_key_data;
    krb5_key_data *key_data;
} osa_pw_hist_ent, *osa_pw_hist_t;

typedef struct _osa_princ_ent_t {
    char * policy;
    u_int32 aux_attributes;
```

```

    unsigned int old_key_len;
    unsigned int old_key_next;
    krb5_kvno admin_history_kvno;
    osa_pw_hist_ent *old_keys;

    u_int32 num_old_keys;
    u_int32 next_old_key;
    krb5_kvno admin_history_kvno;
    osa_pw_hist_ent *old_keys;
} osa_princ_ent_rec, *osa_princ_ent_t;

```

The fields that are different from `kadm5_principal_ent_t` are:

**num\_old\_keys** The number of previous keys in the `old_keys` array. This value must be  $0 \leq \text{num\_old\_keys} < \text{pw\_history\_num}$ .

**old\_key\_next** The index into `old_keys` where the next key should be inserted. This value must be  $0 \leq \text{old\_key\_next} \leq \text{num\_old\_keys}$ .

**admin\_history\_kvno** The key version number of the `kadmin/history` principal's key used to encrypt the values in `old_keys`. If the server library finds that `kadmin/history`'s `kvno` is different from the value in this field, it returns `KADM5_BAD_HIST_KEY`.

**old\_keys** The array of the principal's previous passwords, each encrypted in the `kadmin/history` key. There are `num_old_keys` elements. Each "password" in the array is itself an array of `n_key_data` `krb5_key_data` structures, one for each keysalt type the password was encoded in.

## 6.2 Policy, `osa_policy_ent_t`

The policy database stores records of the type `osa_policy_ent_t` (declared in `<kadm5/adb.h>`), which is all of `kadm5_policy_ent_t` plus necessary bookkeeping information. The records are keyed by the policy name.

```

typedef struct _osa_policy_ent_t {
    char *policy;

    u_int32 pw_min_life;

```

```

    u_int32 pw_max_life;
    u_int32 pw_min_length;
    u_int32 pw_min_classes;
    u_int32 pw_history_num;

    u_int32 refcnt;
} osa_policy_ent_rec, *osa_policy_ent_t;

```

### 6.3 Kerberos, krb5\_db\_entry

The Kerberos database stores records of type `krb5_db_entry`, which is defined in the `<k5-int.h>` header file. The semantics of each field are defined in the `libkdb` functional specification.

## 7 Database Access Methods

### 7.1 Principal and Policy Databases

This section describes the database abstraction used for the admin policy database; the admin principal database used to be treated in the same manner but is now handled more directly as `krb5_tl_data`; thus, nothing in this section applies to it any more. Since both databases export equivalent functionality, the API is only described once. The character `T` is used to represent both “princ” and “policy”. The location of the principal database is defined by the configuration parameters given to any of the `kadm5_init` functions in the server library.

Note that this is *only* a database abstraction. All functional intelligence, such as maintaining policy reference counts or sanity checking, must be implemented above this layer.

Prototypes for the `osa` functions are supplied in `<kadm5/adb.h>`. The routines are defined in `libkadm5srv.a`. They require linking with the Berkely DB library.

#### 7.1.1 Error codes

The database routines use `com_err` for error codes. The error code table name is “adb” and the offsets are the same as the order presented here. The error table header file is `<kadm5/adb_err.h>`. Callers of the OSA routines should first call `init_adb_err_tbl()` to initialize the database table.

**OSA\_ADB\_OK** Operation successful.

**OSA\_ADB\_FAILURE** General failure.

**OSA\_ADB\_DUP** Operation would create a duplicate database entry.

**OSA\_ADB\_NOENT** Named entry not in database.

**OSA\_ADB\_BAD\_PRINC** The krb5\_principal structure is invalid.

**OSA\_ADB\_BAD\_POLICY** The specified policy name is invalid.

**OSA\_ADB\_XDR\_FAILURE** The principal or policy structure cannot be encoded for storage.

**OSA\_ADB\_BADLOCKMODE** Bad lock mode specified.

**OSA\_ADB\_CANTLOCK\_DB** Cannot lock database, presumably because it is already locked.

**OSA\_ADB\_NOTLOCKED** Internal error, database not locked when unlock is called.

**OSA\_ADB\_NOLOCKFILE** KADM5 administration database lock file missing.

Database functions can also return system errors. Unless otherwise specified, database functions return **OSA\_ADB\_OK**.

### 7.1.2 Locking

All of the `osa_adb` functions except `open` and `close` lock and unlock the database to prevent concurrency collisions. The overall locking algorithm is as follows:

1. `osa_adb_open_T` calls `osa_adb_init_db` to allocate the `osa_adb.T.t` structure and open the locking file for further use.
2. Each `osa_adb` functions locks the locking file and opens the appropriate database with `osa_adb_open_and_lock`, performs its action, and then closes the database and unlocks the locking file with `osa_adb_close_and_unlock`.
3. `osa_adb_close_T` calls `osa_adb_fini_db` to close the locking file and deallocate the db structure.

Functions which modify the database acquire an exclusive lock, others acquire a shared lock. `osa_adb_iter_T` acquires an exclusive lock for safety but as stated below consequences of modifying the database in the iteration function are undefined.

### 7.1.3 Function descriptions

```
osa_adb_ret_t osa_adb_create_T_db(kadm5_config_params *params)
```

Create the database and lockfile specified in params. The database must not already exist, or EEXIST is returned. The lock file is only created after the database file has been created successfully.

```
osa_adb_ret_t osa_adb_rename_T_db(kadm5_config_params *fromparams,  
    kadm5_config_params *toparams)
```

Rename the database named by fromparams to that named by toparams. The fromparams database must already exist; the toparams database may exist or not. When the function returns, the database named by fromparams no longer exists, and toparams has been overwritten with fromparams. This function acquires a permanent lock on both databases for the duration of its operation, so a failure is likely to leave the databases unusable.

```
osa_adb_ret_t osa_adb_destroy_policy_db(kadm5_config_params *params)
```

Destroy the database named by params. The database file and lock file are deleted.

```
osa_adb_ret_t  
osa_adb_open_T(osa_adb_T_t *db, char *filename);
```

Open the database named filename. Returns OSA\_ADB\_NOLOCKFILE if the database does not exist or if the lock file is missing. The database is not actually opened in the operating-system file sense until a lock is acquire.

```
osa_adb_ret_t  
osa_adb_close_T(osa_adb_T_t db);
```

Release all shared or exclusive locks (on BOTH databases, since they use the same lock file) and close the database.

It is an error to exit while a permanent lock is held; OSA\_ADB\_NOLOCKFILE is returned in this case.

```
osa_adb_ret_t osa_adb_get_lock(osa_adb_T_t db, int mode)
```

Acquire a lock on the administration databases; note that both databases are locked simultaneously by a single call. The mode argument can be OSA\_ADB\_SHARED, OSA\_ADB\_EXCLUSIVE, or OSA\_ADB\_PERMANENT. The first two and the third are really disjoint locking semantics and should not be interleaved.

Shared and exclusive locks have the usual semantics, and a program can upgrade a shared lock to an exclusive lock by calling the function again. A reference count of open locks is maintained by this function and `osa_adb_release_lock` so the functions can be called multiple times; the actual lock is not released until the final `osa_adb_release_lock`. Note, however, that once a lock is upgraded from shared to exclusive, or from exclusive to permanent, it is not downgraded again until released completely. In other words, `get_lock(SHARED)`, `get_lock(EXCLUSIVE)`, `release_lock()` leaves the process with an exclusive lock with a reference count of one. An attempt to get a shared or exclusive lock that conflicts with another process results in the OSA\_ADB\_CANLOCK\_DB error code.

This function and `osa_adb_release_lock` are called automatically as needed by all other `osa_adb` functions to acquire shared and exclusive locks and so are not normally needed. They can be used explicitly by a program that wants to perform multiple `osa_adb` functions within the context of a single lock.

Acquiring an OSA\_ADB\_PERMANENT lock is different. A permanent lock consists of first acquiring an exclusive lock and then *deleting the lock file*. Any subsequent attempt to acquire a lock by a different process will fail with OSA\_ADB\_NOLOCKFILE instead of OSA\_ADB\_CANLOCK\_DB (attempts in the same process will “succeed” because only the reference count gets incremented). The lock file is recreated by `osa_adb_release_lock` when the last pending lock is released.

The purpose of a permanent lock is to absolutely ensure that the database remain locked during non-atomic operations. If the locking process dies while holding a permanent lock, all subsequent `osa_adb` operations will fail, even through a system reboot. This is useful, for example, for `ovsec_adm_import` which creates both new database files in a temporary location and renames them into place. If both renames do not fully complete the database will probably be inconsistent and everything should stop working until an administrator can clean it up.

```
osa_adb_ret_t osa_adb_release_lock(osa_adb_T_t db)
```

Releases a shared, exclusive, or permanent lock acquired with `osa_adb_get_lock`, or just decrements the reference count if multiple locks are held. When a permanent lock is released, the lock file is re-created.

All of a process' shared or exclusive database locks are released when the process terminates. A permanent lock is *not* released when the process exits (although the exclusive lock it begins

with obviously is).

```
osa_adb_ret_t  
osa_adb_create_T(osa_adb_T_t db, osa_T_ent_t entry);
```

Adds the entry to the database. All fields are defined. Returns OSA\_ADB\_DUP if it already exists.

```
osa_adb_ret_t  
osa_adb_destroy_T(osa_adb_T_t db, osa_T_t name);
```

Removes the named entry from the database. Returns OSA\_ADB\_NOENT if it does not exist.

```
osa_adb_ret_t  
osa_adb_get_T(osa_adb_T_t db, osa_T_t name,  
              osa_princ_ent_t *entry);
```

Looks up the named entry in the db, and returns it in \*entry in allocated storage that must be freed with osa\_adb\_free\_T. Returns OSA\_ADB\_NOENT if name does not exist, OSA\_ADB\_MEM if memory cannot be allocated.

```
osa_adb_ret_t  
osa_adb_put_T(osa_adb_T_t db, osa_T_ent_t entry);
```

Modifies the existing entry named in entry. All fields must be filled in. Returns OSA\_DB\_NOENT if the named entry does not exist. Note that this cannot be used to rename an entry; rename is implemented by deleting the old name and creating the new one (NOT ATOMIC!).

```
void osa_adb_free_T(osa_T_ent_t);
```

Frees the memory associated with an osa\_T\_ent\_t allocated by osa\_adb\_get\_T.

```
typedef osa_adb_ret_t (*osa_adb_iter_T_func)(void *data,  
                                              osa_T_ent_t entry);
```

```
osa_adb_ret_t osa_adb_iter_T(osa_adb_T_t db, osa_adb_iter_T_func func,  
                             void *data);
```

Iterates over every entry in the database. For each entry `ent` in the database `db`, the function `(*func)(data, ent)` is called. If `func` returns an error code, `osa_adb_iter_T` returns an error code. If all invocations of `func` return `OSA_ADB_OK`, `osa_adb_iter_T` returns `OSA_ADB_OK`. The function `func` is permitted to access the database, but the consequences of modifying the database during the iteration are undefined.

## 7.2 Kerberos Database

Kerberos uses the `libkdb` interface to store `krb5_db_entry` records. It can be accessed and modified in parallel with the Kerberos server, using functions that are defined inside the KDC and the `libkdb.a`. The `libkdb` interface is defined in the `libkdb` functional specifications.

### 7.2.1 Initialization and Key Access

Keys stored in the Kerberos database are encrypted in the Kerberos master key. The admin server will therefore have to acquire the key before it can perform any key-changing operations, and will have to decrypt and encrypt the keys retrieved from and placed into the database via `krb5_db_get_principal` and `_put_principal`. This section describes the internal admin server API that will be used to perform these functions.

```
krb5_principal master_princ;  
krb5_encrypt_block master_encblock;  
krb5_keyblock master_keyblock;
```

```
void kdc_init_master()
```

`kdc_init_master` opens the database and acquires the master key. It also sets the global variables `master_princ`, `master_encblock`, and `master_keyblock`:

- `master_princ` is set to the name of the Kerberos master principal (`K/M@REALM`).
- `master_encblock` is something I have no idea about.
- `master_keyblock` is the Kerberos master key

```
krb5_error_code kdb_get_entry_and_key(krb5_principal principal,  
                                     krb5_db_entry *entry,  
                                     krb5_keyblock *key)
```

`kdb_get_entry_and_key` retrieves the named principal's entry from the database in `entry`, and decrypts its key into `key`. The caller must free `entry` with `krb5_dbm_db_free_principal` and free `key->contents` with `free`.<sup>2</sup>

```
krb5_error_code kdb_put_entry_pw(krb5_db_entry *entry, char *pw)
```

`kdb_put_entry_pw` stores `entry` in the database. All the entry values must already be set; this function does not change any of them except the key. `pw`, the NULL-terminated password string, is converted to a key using string-to-key with the salt type specified in `entry->salt_type`.<sup>3</sup>

## 8 Admin Principal and Policy Database Implementation

The admin principal and policy databases will each be stored in a single hash table, implemented by the Berkeley 4.4BSD db library. Each record will consist of an entire `osa_T_ent_t`. The key into the hash table is the entry name (for principals, the ASCII representation of the name). The value is the T entry structure. Since the key and data must be self-contained, with no pointers, the Sun xdr mechanisms will be used to marshal and unmarshal data in the database.

The server in the first release will be single-threaded in that a request will run to completion (or error) before the next will run, but multiple connections will be allowed simultaneously.

## 9 ACLs, `acl_check`

The ACL mechanism described in the “Authorization ACLs” section of the functional specifications will be implemented by the `acl_check` function.

```
enum access_t {  
    ACCESS_DENIED = 0,  
    ACCESS_OK = 1,  
};
```

---

<sup>2</sup>The caller should also `memset(key->contents, 0, key->length)`. There should be a function `krb5_free_keyblock.contents` for this, but there is not.

<sup>3</sup>The `salt_type` should be set based on the command line arguments to the `kadmin` server (see the “Command Line” section of the functional specification).

```
};
```

```
enum access_t acl_check(krb5_principal princ, char *priv);
```

The `priv` argument must be one of “get”, “add”, “delete”, or “modify”. `acl_check` returns 1 if the principal `princ` has the named privilege, 0 if it does not.

## 10 Function Details

This section discusses specific design issues for Admin API functions that are not addressed by the functional specifications.

### 10.1 `kadm5_create_principal`

If the named principal exists in either the Kerberos or admin principal database, but not both, return `KADM5_BAD_DB`.

The principal’s initial key is not stored in the key history array at creation time.

### 10.2 `kadm5_delete_principal`

If the named principal exists in either the Kerberos or admin principal database, but not both, return `KADM5_BAD_DB`.

### 10.3 `kadm5_modify_principal`

If the named principal exists in either the Kerberos or admin principal database, but not both, return `KADM5_BAD_DB`.

If `pw_history_num` changes and the new value  $n$  is smaller than the current value of `num_old_keys`, `old_keys` should end up with the  $n$  most recent keys; these are found by counting backwards  $n$  elements in `old_keys` from `old_key_next`. `old_key_nexts` should then be reset to 0, the oldest of the saved keys, and `num_old_keys` set to  $n$ , the new actual number of old keys in the array.

## 10.4 `kadm5_chpass_principal`, `randkey_principal`

The algorithm for determining whether a password is in the principal's key history is complicated by the use of the `kadmin/history` `K_h` encrypting key.

1. For `kadm5_chpass_principal`, convert the password to a key using `string-to-key` and the salt method specified by the command line arguments.
2. If the `POLICY` bit is set and `pw_history_num` is not zero, check if the new key is in the history.
  - (a) Retrieve the principal's current key and decrypt it with `K_M`. If it is the same as the new key, return `KADM5_PASS_REUSE`.
  - (b) Retrieve the `kadmin/history` key `K_h` and decrypt it with `K_M`.
  - (c) Encrypt the principal's new key in `K_h`.
  - (d) If the principal's new key encrypted in `K_h` is in `old_keys`, return `KADM5_PASS_REUSE`.
  - (e) Encrypt the principal's current key in `K_h` and store it in `old_keys`.
  - (f) Erase the memory containing `K_h`.
3. Encrypt the principal's new key in `K_M` and store it in the database.
4. Erase the memory containing `K_M`.

To store the an encrypted key in `old_keys`, insert it as the `old_key_next` element of `old_keys`, and increment `old_key_next` by one modulo `pw_history_num`.

## 10.5 `kadm5_get_principal`

If the named principal exists in either the Kerberos or admin principal database, but not both, return `KADM5_BAD_DB`.